## Print This Solution Close This Solution

Find solution using dual-simplex method
MAX Z $=\mathbf{- 2 \times 1} \mathbf{- 2 \times 2 - 4 \times 3}$
subject to
$2 \times 1+3 \times 2+5 \times 3>=2$
$3 \times 1+\times 2+7 \times 3<=3$
$\mathrm{x} 1+4 \times 2+6 \times 3<=5$
and $x 1, x 2, x 3>=0$

## Solution:

Problem is
$\operatorname{Max} Z=-2 x_{1}-2 x_{2}-4 x_{3}$
subject to

$$
\begin{aligned}
2 x_{1}+3 x_{2} & +5 x_{3} \geq 2 \\
3 x_{1}+x_{2} & +7 x_{3} \leq 3 \\
x_{1}+4 x_{2} & +6 x_{3} \leq 5 \\
\text { and } x_{1}, x_{2}, x_{3} & \geq 0
\end{aligned}
$$

In order to apply the dual simplex method, convert all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-2 x_{1}-2 x_{2}-4 x_{3}$
subject to

$$
\begin{aligned}
-2 x_{1}-3 x_{2}-5 x_{3} & \leq-2 \\
3 x_{1} & +x_{2}+7 x_{3} \leq 3 \\
x_{1}+4 x_{2} & +6 x_{3} \leq 5
\end{aligned}
$$

and $x_{1}, x_{2}, x_{3} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$
3. As the constraint 3 is of type ' $\leq$ ' we should add slack variable $S_{3}$

## After introducing slack variables

$\operatorname{Max} Z=-2 x_{1}-2 x_{2}-4 x_{3}+0 S_{1}+0 S_{2}+0 S_{3}$
subject to

$$
\begin{aligned}
-2 x_{1}-3 x_{2}-5 x_{3}+S_{1} & =-2 \\
3 x_{1}+x_{2}+7 x_{3}+S_{2} & =3 \\
x_{1}+4 x_{2}+6 x_{3} & +S_{3}
\end{aligned}=5
$$

and $x_{1}, x_{2}, x_{3}, S_{1}, S_{2}, S_{3} \geq 0$

| Iteration-1 |  | $C_{j}$ | -2 | -2 | -4 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $S_{1}$ | $S_{2}$ | $S_{3}$ |
| $S_{1}$ | 0 | -2 | -2 | (-3) | -5 | 1 | 0 | 0 |
| $S_{2}$ | 0 | 3 | 3 | 1 | 7 | 0 | 1 | 0 |
| $S_{3}$ | 0 | 5 | 1 | 4 | 6 | 0 | 0 | 1 |
| $Z=0$ |  | $Z_{j}$ | 0 | 0 | 0 | 0 | 0 | 0 |
|  |  | $C_{j}-Z_{j}$ | -2 | -2 | -4 | 0 | 0 | 0 |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{1}, j} \\ \text { and } S_{1}, j<0 \end{array}$ | 1 | $\frac{2}{3} \uparrow$ | $\frac{4}{5}$ | --- | --- | --- |

Minimum negative $X_{B}$ is -2 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is $\frac{2}{3}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is -3 .
Entering $=x_{2}$, Departing $=S_{1}$, Key Element $=-3$
$R_{1}($ new $)=R_{1}($ old $) \div-3$
$R_{2}$ (new) $=R_{2}($ old $)-R_{1}$ (new)
$R_{3}$ (new) $=R_{3}$ (old) $-4 R_{1}$ (new)

| Iteration-2  $C_{j}$ -2 -2 -4 0 0 0 <br> $\boldsymbol{B}$ $\boldsymbol{C}_{\boldsymbol{B}}$ $\boldsymbol{X}_{\boldsymbol{B}}$ $\boldsymbol{x}_{\mathbf{1}}$ $\boldsymbol{x}_{\mathbf{2}}$ $\boldsymbol{x}_{\mathbf{3}}$ $\boldsymbol{S}_{\mathbf{1}}$ $\boldsymbol{S}_{\mathbf{2}}$ $\boldsymbol{S}_{\mathbf{3}}$ <br> $x_{2}$ -2 $\frac{2}{3}$ $\frac{2}{3}$ 1 $\frac{5}{3}$ $-\frac{1}{3}$ 0 0 <br> $S_{2}$ 0 $\frac{7}{3}$ $\frac{7}{3}$ 0 $\frac{16}{3}$ $\frac{1}{3}$ 1 0 <br> $S_{3}$ 0 $\frac{7}{3}$ $-\frac{5}{3}$ 0 $-\frac{2}{3}$ $\frac{4}{3}$ 0 1 <br> $\boldsymbol{Z}=-\frac{\mathbf{4}}{\mathbf{3}}$  $\boldsymbol{Z}_{\boldsymbol{j}}$ $-\frac{\mathbf{4}}{\mathbf{3}}$ $\mathbf{- 2}$ $\mathbf{- \frac { 1 } { 3 }}$ $\frac{\mathbf{2}}{\mathbf{3}}$ $\mathbf{0}$ $\mathbf{0}$ |
| :--- |


|  |  | $C_{j}-Z_{j}$ | $-\frac{2}{3}$ | 0 | $-\frac{2}{3}$ | $-\frac{2}{3}$ | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Ratio | --- | --- | --- | --- | --- | --- |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=0, x_{2}=\frac{2}{3}, x_{3}=0$
$\operatorname{Max} Z=-\frac{4}{3}$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MAX Z $=-\mathrm{x} 1-\mathbf{2 x} \mathbf{2}$
subject to
$-2 \times 1-\times 2<=-4$
$-\mathrm{x} 1-2 \times 2<=-7$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Max} Z=-x_{1}-2 x_{2}$
subject to
$-2 x_{1}-x_{2} \leq-4$

- $x_{1}-2 x_{2} \leq-7$
and $x_{1}, x_{2} \geq 0 ;$

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-x_{1}-2 x_{2}+0 S_{1}+0 S_{2}$
subject to
$-2 x_{1}-x_{2}+S_{1}=-4$
$-x_{1}-2 x_{2}+S_{2}=-7$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -1 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | -4 | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ |
| $S_{1}$ | 0 | -7 | $\boldsymbol{S}_{\mathbf{2}}$ |  |  |  |
| $\boldsymbol{S}_{\mathbf{2}}$ | 0 | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $\mathbf{( - 1 )}$ | -2 | 0 | 1 |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $C_{j}-Z_{j}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  | Ratio $=\frac{C_{j}-Z_{j}}{S_{2}, j}$ <br> and $S_{2}, j<0$ | -1 | -2 | 0 | 0 |  |
|  |  | 1 | 1 | --- | --- |  |

Minimum negative $X_{B}$ is -7 and its row index is 2 . So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is 1 and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -1 .
Entering $=x_{1}$, Departing $=S_{2}$, Key Element $=-1$
$R_{2}($ new $)=R_{2}($ old $) \div-1$
$R_{1}$ (new) $=R_{1}($ old $)+2 R_{2}$ (new)

| Iteration-2 |  | $C_{j}$ | -1 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | 10 | 0 | 3 | 1 | -2 |
| $x_{1}$ | -1 | 7 | 1 | 2 | 0 | -1 |
| $\boldsymbol{Z}=-7$ |  | $Z_{j}$ | $-\mathbf{1}$ | $-\mathbf{2}$ | $\mathbf{0}$ | $\boldsymbol{1}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | 0 | 0 | -1 |
|  | Ratio | --- | --- | $\boldsymbol{- -}$ | $\boldsymbol{- -}$ |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=7, x_{2}=0$
$\operatorname{Max} Z=-7$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=2 \times 1+\mathbf{x} 2$
subject to
$3 \times 2>=6$
$3 \times 1+2 \times 2>=8$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=2 x_{1}+x_{2}$
subject to

$$
3 x_{2} \geq 6
$$

$3 x_{1}+2 x_{2} \geq 8$
and $x_{1}, x_{2} \geq 0$;

In order to apply the dual simplex method, convert Min $Z$ to Max $Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-2 x_{1}-x_{2}$
subject to
$-3 x_{2} \leq-6$
$-3 x_{1}-2 x_{2} \leq-8$
and $x_{1}, x_{2} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-2 x_{1}-x_{2}+0 S_{1}+0 S_{2}$
subject to

$$
-3 x_{2}+S_{1}=-6
$$

$-3 x_{1}-2 x_{2}+S_{2}=-8$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -2 | -1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | -6 | 0 | -3 | 1 | 0 |
|  |  |  |  |  |  |  |
| about:blank |  |  |  |  |  |  |


| $\boldsymbol{S}_{\mathbf{2}}$ | 0 | -8 | -3 | $\mathbf{( - 2 )}$ | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{Z = 0}$ |  | $Z_{\boldsymbol{j}}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  |  | $C_{j}-Z_{j}$ | -2 | -1 | 0 | 0 |
|  | Ratio $=\frac{C_{j}-Z_{j}}{S_{2}, j}$ <br> and $S_{2}, j<0$ | $\frac{2}{3}$ | $\frac{1}{2} \uparrow$ | --- | --e |  |

Minimum negative $X_{B}$ is -8 and its row index is 2 . So, the leaving basis variable is $S_{2}$.

Minimum positive ratio is $\frac{1}{2}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is -2 .
Entering $=x_{2}$, Departing $=S_{2}$, Key Element $=-2$
$R_{2}($ new $)=R_{2}($ old $) \div-2$
$R_{1}($ new $)=R_{1}($ old $)+3 R_{2}($ new $)$

| Iteration-2 |  | $C_{j}$ | -2 | -1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | 6 | $\frac{9}{2}$ | 0 | 1 | $-\frac{3}{2}$ |
| $x_{2}$ | -1 | 4 | $\frac{3}{2}$ | 1 | 0 | $-\frac{1}{2}$ |
| $Z=-4$ |  | $Z_{j}$ | $-\frac{\mathbf{3}}{\mathbf{2}}$ | $-\mathbf{1}$ | $\mathbf{0}$ | $\frac{\mathbf{1}}{\mathbf{2}}$ |
|  |  | $C_{j}-Z_{j}$ | $-\frac{1}{2}$ | 0 | 0 | $-\frac{1}{2}$ |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=0, x_{2}=4$
$\operatorname{Max} Z=-4$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN Z $=3 \times 1+4 \times 2$
subject to
$2 \times 1+3 \times 2>=90$
$4 \times 1+3 \times 2>=120$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=3 x_{1}+4 x_{2}$
subject to

$$
\begin{aligned}
& 2 x_{1}+3 x_{2} \geq 90 \\
& 4 x_{1}+3 x_{2} \geq 120 \\
& \text { and } x_{1}, x_{2} \geq 0
\end{aligned}
$$

In order to apply the dual simplex method, convert $\operatorname{Min} Z$ to $\operatorname{Max} Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-3 x_{1}-4 x_{2}$
subject to
$-2 x_{1}-3 x_{2} \leq-90$
$-4 x_{1}-3 x_{2} \leq-120$
and $x_{1}, x_{2} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-3 x_{1}-4 x_{2}+0 S_{1}+0 S_{2}$
subject to
$-2 x_{1}-3 x_{2}+S_{1}=-90$
$-4 x_{1}-3 x_{2}+S_{2}=-120$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -3 | -4 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | -90 | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ |
| $S_{1}$ | 0 | -2 | -3 | $\boldsymbol{S}_{\mathbf{2}}$ |  |  |
|  |  |  |  |  |  |  |
| about:blank |  |  |  |  |  |  |


| $S_{2}$ | 0 | -120 | (-4) | -3 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z=0$ |  | $Z_{j}$ | 0 | 0 | 0 | 0 |
|  |  | $C_{j}-Z_{j}$ | -3 | -4 | 0 | 0 |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{2}, j} \\ \text { and } S_{2}, j<0 \end{array}$ | $\frac{3}{4} \uparrow$ | $\frac{4}{3}$ | --- | --- |

Minimum negative $X_{B}$ is -120 and its row index is 2 . So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is $\frac{3}{4}$ and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -4 .
Entering $=x_{1}$, Departing $=S_{2}$, Key Element $=-4$
$R_{2}($ new $)=R_{2}($ old $) \div-4$
$R_{1}$ (new) $=R_{1}($ old $)+2 R_{2}$ (new)

| Iteration-2 |  | $C_{j}$ | -3 | -4 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $S_{1}$ | $S_{2}$ |
| $S_{1}$ | 0 | -30 | 0 | $\left(-\frac{3}{2}\right)$ | 1 | $-\frac{1}{2}$ |
| $x_{1}$ | -3 | 30 | 1 | $\frac{3}{4}$ | 0 | $-\frac{1}{4}$ |
| $Z=-90$ |  | $Z_{j}$ | -3 | $-\frac{9}{4}$ | 0 | $\frac{3}{4}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | $-\frac{7}{4}$ | 0 | $-\frac{3}{4}$ |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{1}, j} \\ \text { and } S_{1}, j<0 \end{array}$ | --- | $\frac{7}{6} \uparrow$ | --- | $\frac{3}{2}$ |

Minimum negative $X_{B}$ is -30 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is $\frac{7}{6}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is $-\frac{3}{2}$.
Entering $=x_{2}$, Departing $=S_{1}$, Key Element $=-\frac{3}{2}$
$R_{1}($ new $)=R_{1}(\mathrm{old}) \times-\frac{2}{3}$
$R_{2}($ new $)=R_{2}($ old $)-\frac{3}{4} R_{1}$ (new)

| Iteration-3 |  | $C_{j}$ | -3 | -4 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $x_{2}$ | -4 | 20 | 0 | 1 | $-\frac{2}{3}$ | $\frac{1}{3}$ |
| $x_{1}$ | -3 | 15 | 1 | 0 | $\frac{1}{2}$ | $-\frac{1}{2}$ |
| $\boldsymbol{Z}=-\mathbf{1 2 5}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | -3 | -4 | $\frac{\mathbf{7}}{\mathbf{6}}$ | $\frac{\mathbf{1}}{\mathbf{6}}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | 0 | $-\frac{7}{6}$ | $-\frac{1}{6}$ |
|  | Ratio | --- | --- | --- | --- |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=15, x_{2}=20$
$\operatorname{Max} Z=-125$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
$\operatorname{MIN} Z=2 \times 1+2 \times 2+4 \times 3$
subject to
$2 \times 1+3 \times 2+5 \times 3>=2$
$3 \times 1+\mathrm{x} 2+7 \times 3<=3$
$\mathrm{x} 1+4 \times 2+6 \times 3<=5$
and $\mathrm{x} 1, \mathrm{x} 2, \mathrm{x} 3>=0$

## Solution:

Problem is
$\operatorname{Min} Z=2 x_{1}+2 x_{2}+4 x_{3}$
subject to

$$
\begin{aligned}
2 x_{1}+3 x_{2} & +5 x_{3} \geq 2 \\
3 x_{1}+x_{2} & +7 x_{3} \leq 3 \\
x_{1}+4 x_{2} & +6 x_{3} \leq 5 \\
\text { and } x_{1}, x_{2}, x_{3} & \geq 0
\end{aligned}
$$

In order to apply the dual simplex method, convert Min $Z$ to $\operatorname{Max} Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-2 x_{1}-2 x_{2}-4 x_{3}$
subject to

$$
\begin{aligned}
-2 x_{1}-3 x_{2}-5 x_{3} & \leq-2 \\
3 x_{1} & +x_{2}+7 x_{3} \leq 3 \\
x_{1}+4 x_{2} & +6 x_{3} \leq 5
\end{aligned}
$$

and $x_{1}, x_{2}, x_{3} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$
3. As the constraint 3 is of type ' $\leq$ ' we should add slack variable $S_{3}$

## After introducing slack variables

$\operatorname{Max} Z=-2 x_{1}-2 x_{2}-4 x_{3}+0 S_{1}+0 S_{2}+0 S_{3}$
subject to

$$
\begin{aligned}
-2 x_{1}-3 x_{2}-5 x_{3}+S_{1} & =-2 \\
3 x_{1}+x_{2}+7 x_{3}+S_{2} & =3 \\
x_{1}+4 x_{2}+6 x_{3} & +S_{3}
\end{aligned}=5
$$

and $x_{1}, x_{2}, x_{3}, S_{1}, S_{2}, S_{3} \geq 0$

| Iteration-1 |  | $C_{j}$ | -2 | -2 | -4 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $S_{1}$ | $S_{2}$ | $S_{3}$ |
| $S_{1}$ | 0 | -2 | -2 | (-3) | -5 | 1 | 0 | 0 |
| $S_{2}$ | 0 | 3 | 3 | 1 | 7 | 0 | 1 | 0 |
| $S_{3}$ | 0 | 5 | 1 | 4 | 6 | 0 | 0 | 1 |
| $Z=0$ |  | $Z_{j}$ | 0 | 0 | 0 | 0 | 0 | 0 |
|  |  | $C_{j}-Z_{j}$ | -2 | -2 | -4 | 0 | 0 | 0 |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{1}, j} \\ \text { and } S_{1}, j<0 \end{array}$ | 1 | $\frac{2}{3} \uparrow$ | $\frac{4}{5}$ | --- | --- | --- |

Minimum negative $X_{B}$ is -2 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is $\frac{2}{3}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is -3 .
Entering $=x_{2}$, Departing $=S_{1}$, Key Element $=-3$
$R_{1}($ new $)=R_{1}($ old $) \div-3$
$R_{2}$ (new) $=R_{2}($ old $)-R_{1}$ (new)
$R_{3}$ (new) $=R_{3}$ (old) $-4 R_{1}$ (new)

| Iteration-2  $C_{j}$ -2 -2 -4 0 0 0 <br> $\boldsymbol{B}$ $\boldsymbol{C}_{\boldsymbol{B}}$ $\boldsymbol{X}_{\boldsymbol{B}}$ $\boldsymbol{x}_{\mathbf{1}}$ $\boldsymbol{x}_{\mathbf{2}}$ $\boldsymbol{x}_{\mathbf{3}}$ $\boldsymbol{S}_{\mathbf{1}}$ $\boldsymbol{S}_{\mathbf{2}}$ $\boldsymbol{S}_{\mathbf{3}}$ <br> $x_{2}$ -2 $\frac{2}{3}$ $\frac{2}{3}$ 1 $\frac{5}{3}$ $-\frac{1}{3}$ 0 0 <br> $S_{2}$ 0 $\frac{7}{3}$ $\frac{7}{3}$ 0 $\frac{16}{3}$ $\frac{1}{3}$ 1 0 <br> $S_{3}$ 0 $\frac{7}{3}$ $-\frac{5}{3}$ 0 $-\frac{2}{3}$ $\frac{4}{3}$ 0 1 <br> $\boldsymbol{Z}=-\frac{\mathbf{4}}{\mathbf{3}}$  $\boldsymbol{Z}_{\boldsymbol{j}}$ $-\frac{\mathbf{4}}{\mathbf{3}}$ $\mathbf{- 2}$ $\mathbf{- \frac { 1 } { 3 }}$ $\frac{\mathbf{2}}{\mathbf{3}}$ $\mathbf{0}$ $\mathbf{0}$ |
| :--- |


|  |  | $C_{j}-Z_{j}$ | $-\frac{2}{3}$ | 0 | $-\frac{2}{3}$ | $-\frac{2}{3}$ | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Ratio | --- | --- | --- | --- | --- | --- |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=0, x_{2}=\frac{2}{3}, x_{3}=0$
$\operatorname{Max} Z=-\frac{4}{3}$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=2 \times 1+\times 2+4 \times 3$
subject to
$2 \times 1+3 \times 2+3 \times 3>=12$
$3 \times 1+2 \times 2+\times 3>=18$
and $\mathrm{x} 1, \mathrm{x} 2, \mathrm{x} 3>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=2 x_{1}+x_{2}+4 x_{3}$
subject to

$$
\begin{aligned}
& 2 x_{1}+3 x_{2}+3 x_{3} \geq 12 \\
& 3 x_{1}+2 x_{2}+x_{3} \geq 18 \\
& \text { and } x_{1}, x_{2}, x_{3} \geq 0 ;
\end{aligned}
$$

In order to apply the dual simplex method, convert $\operatorname{Min} Z$ to $\operatorname{Max} Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-2 x_{1}-x_{2}-4 x_{3}$
subject to
$-2 x_{1}-3 x_{2}-3 x_{3} \leq-12$
$-3 x_{1}-2 x_{2}-x_{3} \leq-18$
and $x_{1}, x_{2}, x_{3} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-2 x_{1}-x_{2}-4 x_{3}+0 S_{1}+0 S_{2}$
subject to
$-2 x_{1}-3 x_{2}-3 x_{3}+S_{1}=-12$
$-3 x_{1}-2 x_{2}-x_{3}+S_{2}=-18$
and $x_{1}, x_{2}, x_{3}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -2 | -1 | -4 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{x}_{\mathbf{3}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | -12 | -2 | -3 | -3 | 1 | 0 |
|  |  |  |  |  |  |  |  |
| about:blank |  |  |  |  |  |  |  |


| $S_{2}$ | 0 | -18 | -3 | (-2) | -1 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z=0$ |  | $Z_{j}$ | 0 | 0 | 0 | 0 | 0 |
|  |  | $C_{j}-Z_{j}$ | -2 | -1 | -4 | 0 | 0 |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{2}, j} \\ \text { and } S_{2}, j<0 \end{array}$ | $\frac{2}{3}$ | $\frac{1}{2} \uparrow$ | 4 | -- | --- |

Minimum negative $X_{B}$ is -18 and its row index is 2 . So, the leaving basis variable is $S_{2}$.

Minimum positive ratio is $\frac{1}{2}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is -2 .
Entering $=x_{2}$, Departing $=S_{2}$, Key Element $=-2$
$R_{2}($ new $)=R_{2}($ old $) \div-2$
$R_{1}($ new $)=R_{1}($ old $)+3 R_{2}($ new $)$

| Iteration-2 |  | $C_{j}$ | -2 | -1 | -4 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{x}_{\mathbf{3}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | 15 | $\frac{5}{2}$ | 0 | $-\frac{3}{2}$ | 1 | $-\frac{3}{2}$ |
| $x_{2}$ | -1 | 9 | $\frac{3}{2}$ | 1 | $\frac{1}{2}$ | 0 | $-\frac{1}{2}$ |
| $\boldsymbol{Z}=-\mathbf{9}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $-\frac{\mathbf{3}}{\mathbf{2}}$ | $-\mathbf{1}$ | $-\frac{\mathbf{1}}{\mathbf{2}}$ | $\mathbf{0}$ | $\frac{\mathbf{1}}{\mathbf{2}}$ |
|  |  | $C_{j}-Z_{j}$ | $-\frac{1}{2}$ | 0 | $-\frac{7}{2}$ | 0 | $-\frac{1}{2}$ |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=0, x_{2}=9, x_{3}=0$
$\operatorname{Max} Z=-9$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=3 \times 1+2 \times 2+x 3$
subject to
$2 \times 1+\mathrm{x} 2+4 \times 3>=15$
$\mathrm{x} 1+4 \times 2+3 \times 3>=21$
and $\mathrm{x} 1, \mathrm{x} 2, \mathrm{x} 3>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=3 x_{1}+2 x_{2}+x_{3}$
subject to

$$
\begin{aligned}
2 x_{1}+x_{2}+4 x_{3} & \geq 15 \\
x_{1}+4 x_{2}+3 x_{3} & \geq 21
\end{aligned}
$$

and $x_{1}, x_{2}, x_{3} \geq 0$;

In order to apply the dual simplex method, convert $\operatorname{Min} Z$ to $\operatorname{Max} Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-3 x_{1}-2 x_{2}-x_{3}$
subject to
$-2 x_{1}-x_{2}-4 x_{3} \leq-15$

- $x_{1}-4 x_{2}-3 x_{3} \leq-21$
and $x_{1}, x_{2}, x_{3} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-3 x_{1}-2 x_{2}-x_{3}+0 S_{1}+0 S_{2}$
subject to
$-2 x_{1}-x_{2}-4 x_{3}+S_{1}=-15$

- $x_{1}-4 x_{2}-3 x_{3}+S_{2}=-21$
and $x_{1}, x_{2}, x_{3}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -3 | -2 | -1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{x}_{\mathbf{3}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | -15 | -2 | -1 | -4 | 1 | 0 |


| $S_{2}$ | 0 | -21 | -1 | -4 | (-3) | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z=0$ |  | $Z_{j}$ | 0 | 0 | 0 | 0 | 0 |
|  |  | $C_{j}-Z_{j}$ | -3 | -2 | -1 | 0 | 0 |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{2}, j} \\ \text { and } S_{2}, j<0 \end{array}$ | 3 | $\frac{1}{2}$ | $\frac{1}{3} \uparrow$ | - | --- |

Minimum negative $X_{B}$ is -21 and its row index is 2 . So, the leaving basis variable is $S_{2}$.

Minimum positive ratio is $\frac{1}{3}$ and its column index is 3 . So, the entering variable is $x_{3}$.
$\therefore$ The pivot element is -3 .
Entering $=x_{3}$, Departing $=S_{2}$, Key Element $=-3$
$R_{2}($ new $)=R_{2}($ old $) \div-3$
$R_{1}($ new $)=R_{1}($ old $)+4 R_{2}($ new $)$

| Iteration-2 |  | $C_{j}$ | -3 | -2 | -1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{x}_{\mathbf{3}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | 13 | $-\frac{2}{3}$ | $\frac{13}{3}$ | 0 | 1 | $-\frac{4}{3}$ |
| $x_{3}$ | -1 | 7 | $\frac{1}{3}$ | $\frac{4}{3}$ | 1 | 0 | $-\frac{1}{3}$ |
| $\boldsymbol{Z}=-7$ | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $-\frac{\mathbf{1}}{\mathbf{3}}$ | $-\frac{4}{\mathbf{3}}$ | $-\mathbf{1}$ | $\mathbf{0}$ | $\frac{\mathbf{1}}{\mathbf{3}}$ |  |
|  | $C_{j}-Z_{j}$ | $-\frac{8}{3}$ | $-\frac{2}{3}$ | 0 | 0 | $-\frac{1}{3}$ |  |
|  | Ratio | --- | --- | --- | --- | --- |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=0, x_{2}=0, x_{3}=7$
$\operatorname{Max} Z=-7$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
$\operatorname{MIN} Z=3 \times 1+\mathbf{x} 2+2 \times 3$
subject to
$4 \times 1+\mathrm{x} 2+4 \times 3>=12$
$\mathrm{x} 1+3 \times 2+4 \times 3>=10$
$2 \times 1+2 \times 2+\mathrm{x} 3>=6$
and $\mathrm{x} 1, \mathrm{x} 2, \mathrm{x} 3>=0$

## Solution:

Problem is
$\operatorname{Min} Z=3 x_{1}+x_{2}+2 x_{3}$
subject to

$$
\begin{aligned}
4 x_{1}+x_{2}+4 x_{3} & \geq 12 \\
x_{1}+3 x_{2}+4 x_{3} & \geq 10 \\
2 x_{1}+2 x_{2}+x_{3} & \geq 6 \\
\text { and } x_{1}, x_{2}, x_{3} & \geq 0
\end{aligned}
$$

In order to apply the dual simplex method, convert Min $Z$ to $\operatorname{Max} Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-3 x_{1}-x_{2}-2 x_{3}$
subject to
$-4 x_{1}-x_{2}-4 x_{3} \leq-12$

- $x_{1}-3 x_{2}-4 x_{3} \leq-10$
$-2 x_{1}-2 x_{2}-x_{3} \leq-6$
and $x_{1}, x_{2}, x_{3} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$
3. As the constraint 3 is of type ' $\leq$ ' we should add slack variable $S_{3}$

## After introducing slack variables

$\operatorname{Max} Z=-3 x_{1}-x_{2}-2 x_{3}+0 S_{1}+0 S_{2}+0 S_{3}$
subject to
$-4 x_{1}-x_{2}-4 x_{3}+S_{1} \quad=-12$
$-x_{1}-3 x_{2}-4 x_{3}+S_{2}=-10$
$-2 x_{1}-2 x_{2}-x_{3}+S_{3}=-6$
and $x_{1}, x_{2}, x_{3}, S_{1}, S_{2}, S_{3} \geq 0$

| Iteration-1 |  | $C_{j}$ | -3 | -1 | -2 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | -12 | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{x}_{\mathbf{3}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $\boldsymbol{S}_{\mathbf{1}}$ | 0 | -10 | -4 | -1 | $\mathbf{( - 4 )}$ | 1 | 0 | 0 |
| $S_{2}$ | 0 | -6 | -1 | -3 | -4 | 0 | 1 | 0 |
| $S_{3}$ | 0 | $\boldsymbol{Z}_{\boldsymbol{j}}$ | -2 | -2 | -1 | 0 | 0 | 1 |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $C_{j}-Z_{j}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  | $C_{j}-Z_{j}$ <br> Ratio $=\frac{S_{1}, j}{}$ <br> and $S_{1}, j<0$ | $\frac{3}{4}$ | 1 | $\frac{1}{2} \uparrow$ | --- | --- | --- |  |

Minimum negative $X_{B}$ is -12 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is $\frac{1}{2}$ and its column index is 3 . So, the entering variable is $x_{3}$.
$\therefore$ The pivot element is -4 .
Entering $=x_{3}$, Departing $=S_{1}$, Key Element $=-4$
$R_{1}($ new $)=R_{1}($ old $) \div-4$
$R_{2}($ new $)=R_{2}($ old $)+4 R_{1}$ (new)
$R_{3}($ new $)=R_{3}$ (old) $+R_{1}$ (new)

| Iteration-2 |  | $C_{j}$ | -3 | -1 | -2 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $S_{1}$ | $S_{2}$ | $S_{3}$ |
| $x_{3}$ | -2 | 3 | 1 | $\frac{1}{4}$ | 1 | $-\frac{1}{4}$ | 0 | 0 |
| $S_{2}$ | 0 | 2 | 3 | -2 | 0 | -1 | 1 | 0 |
| $S_{3}$ | 0 | -3 | -1 | $\left(-\frac{7}{4}\right)$ | 0 | - $\frac{1}{4}$ | 0 | 1 |
| $Z=-6$ |  | $Z_{j}$ | -2 | $-\frac{1}{2}$ | -2 | $\frac{1}{2}$ | 0 | 0 |


|  | $C_{j}-Z_{j}$ | -1 | $-\frac{1}{2}$ | 0 | $-\frac{1}{2}$ | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Ratio $=\frac{C_{j}-Z_{j}}{S_{3}, j}$ <br> and $S_{3}, j<0$ | 1 | $\frac{2}{7} \uparrow$ | --- | 2 | --- | --- |
|  |  |  |  |  |  |  |  |

Minimum negative $X_{B}$ is -3 and its row index is 3 . So, the leaving basis variable is $S_{3}$.

Minimum positive ratio is $\frac{2}{7}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is $-\frac{7}{4}$.

Entering $=x_{2}$, Departing $=S_{3}$, Key Element $=-\frac{7}{4}$
$R_{3}($ new $)=R_{3}($ old $) \times-\frac{4}{7}$
$R_{1}($ new $)=R_{1}($ old $)-\frac{1}{4} R_{3}($ new $)$
$R_{2}($ new $)=R_{2}($ old $)+2 R_{3}($ new $)$

| Iteration-3 |  | $C_{j}$ | -3 | -1 | -2 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $\boldsymbol{x}_{3}$ | $S_{1}$ | $S_{2}$ | $S_{3}$ |
| $x_{3}$ | -2 | $\frac{18}{7}$ | $\frac{6}{7}$ | 0 | 1 | $-\frac{2}{7}$ | 0 | $\frac{1}{7}$ |
| $S_{2}$ | 0 | $\frac{38}{7}$ | $\frac{29}{7}$ | 0 | 0 | $-\frac{5}{7}$ | 1 | $-\frac{8}{7}$ |
| $x_{2}$ | -1 | $\frac{12}{7}$ | $\frac{4}{7}$ | 1 | 0 | $\frac{1}{7}$ | 0 | $-\frac{4}{7}$ |
| $Z=-\frac{48}{7}$ |  | $Z_{j}$ | $-\frac{16}{7}$ | -1 | -2 | $\frac{3}{7}$ | 0 | $\frac{2}{7}$ |
|  |  | $C_{j}-Z_{j}$ | $-\frac{5}{7}$ | 0 | 0 | $-\frac{3}{7}$ | 0 | $-\frac{2}{7}$ |
|  |  | Ratio | --- | --- | --- | --- | --- | -- |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.

Hence, optimal solution is arrived with value of variables as :
$x_{1}=0, x_{2}=\frac{12}{7}, x_{3}=\frac{18}{7}$
$\operatorname{Max} Z=-\frac{48}{7}$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=4 \times 1+2 \times 2$
subject to
$4 \times 1+\times 2>=14$
$\mathrm{x} 1+3 \times 2>=12$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=4 x_{1}+2 x_{2}$
subject to

$$
\begin{array}{r}
4 x_{1}+x_{2} \geq 14 \\
x_{1}+3 x_{2} \geq 12
\end{array}
$$

and $x_{1}, x_{2} \geq 0 ;$

In order to apply the dual simplex method, convert $\operatorname{Min} \mathrm{Z}$ to $\operatorname{Max} \mathrm{Z}$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-4 x_{1}-2 x_{2}$
subject to
$-4 x_{1}-x_{2} \leq-14$

- $x_{1}-3 x_{2} \leq-12$
and $x_{1}, x_{2} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-4 x_{1}-2 x_{2}+0 S_{1}+0 S_{2}$
subject to
$-4 x_{1}-x_{2}+S_{1}=-14$

- $x_{1}-3 x_{2}+S_{2}=-12$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -4 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $\boldsymbol{S}_{\mathbf{1}}$ | 0 | -14 | $(-4)$ | -1 | 1 | 0 |
|  |  |  |  |  |  |  |
| about:blank |  |  |  |  |  |  |


| $S_{2}$ | 0 | -12 | -1 | -3 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $Z_{\boldsymbol{j}}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  |  | $C_{j}-Z_{j}$ | -4 | -2 | 0 | 0 |
|  | Ratio $=\frac{C_{j}-Z_{j}}{S_{1}, j}$ <br> and $S_{1}, j<0$ | $1 \uparrow$ | 2 | --- | --- |  |

Minimum negative $X_{B}$ is -14 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is 1 and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -4 .
Entering $=x_{1}$, Departing $=S_{1}$, Key Element $=-4$
$R_{1}($ new $)=R_{1}($ old $) \div-4$
$R_{2}$ (new) $=R_{2}$ (old) $+R_{1}$ (new)

| Iteration-2 |  | $C_{j}$ | -4 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $S_{1}$ | $S_{2}$ |
| $x_{1}$ | -4 | $\frac{7}{2}$ | 1 | $\frac{1}{4}$ | - $\frac{1}{4}$ | 0 |
| $S_{2}$ | 0 | $-\frac{17}{2}$ | 0 | $\left(-\frac{11}{4}\right)$ | - $\frac{1}{4}$ | 1 |
| $Z=-14$ |  | $Z_{j}$ | -4 | -1 | 1 | 0 |
|  |  | $C_{j}-Z_{j}$ | 0 | -1 | -1 | 0 |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{2}, j} \\ \text { and } S_{2}, j<0 \end{array}$ | --- | $\frac{4}{11} \uparrow$ | 4 | --- |

Minimum negative $X_{B}$ is $-\frac{17}{2}$ and its row index is 2. So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is $\frac{4}{11}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is $-\frac{11}{4}$.
Entering $=x_{2}$, Departing $=S_{2}$, Key Element $=-\frac{11}{4}$
$R_{2}($ new $)=R_{2}($ old $) \times-\frac{4}{11}$
$R_{1}$ (new) $=R_{1}($ old $)-\frac{1}{4} R_{2}$ (new)

| Iteration-3 |  | $C_{j}$ | -4 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $x_{1}$ | -4 | $\frac{30}{11}$ | 1 | 0 | $-\frac{3}{11}$ | $\frac{1}{11}$ |
| $x_{2}$ | -2 | $\frac{34}{11}$ | 0 | 1 | $\frac{1}{11}$ | $-\frac{4}{11}$ |
| $\boldsymbol{Z}=-\frac{\mathbf{1 8 8}}{\mathbf{1 1}}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | -4 | $-\mathbf{2}$ | $\frac{\mathbf{1 0}}{\mathbf{1 1}}$ | $\frac{\mathbf{4}}{\mathbf{1 1}}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | 0 | $-\frac{10}{11}$ | $-\frac{4}{11}$ |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=\frac{30}{11}, x_{2}=\frac{34}{11}$
$\operatorname{Max} Z=-\frac{188}{11}$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=5 \times 1+8 \times 2$
subject to
$2 \times 1+3 \times 2>=4$
$\mathrm{x} 1-2 \mathrm{x} 2>=5$
$\mathrm{x} 1+\mathrm{x} 2>=12$
and $x 1, x 2>=0$

## Solution:

Problem is
$\operatorname{Min} Z=5 x_{1}+8 x_{2}$
subject to

$$
\begin{array}{r}
2 x_{1}+3 x_{2} \geq 4 \\
x_{1}-2 x_{2} \geq 5 \\
x_{1}+x_{2} \geq 12 \\
\text { and } x_{1}, x_{2} \geq 0
\end{array}
$$

In order to apply the dual simplex method, convert $\operatorname{Min} \mathrm{Z}$ to $\operatorname{Max} \mathrm{Z}$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-5 x_{1}-8 x_{2}$
subject to
$-2 x_{1}-3 x_{2} \leq-4$
$-x_{1}+2 x_{2} \leq-5$

- $x_{1}-x_{2} \leq-12$
and $x_{1}, x_{2} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ${ }^{\prime} \leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$
3. As the constraint 3 is of type ' $\leq$ ' we should add slack variable $S_{3}$

## After introducing slack variables

$\operatorname{Max} Z=-5 x_{1}-8 x_{2}+0 S_{1}+0 S_{2}+0 S_{3}$
subject to
$-2 x_{1}-3 x_{2}+S_{1} \quad=-4$
$-x_{1}+2 x_{2}+S_{2}=-5$

- $x_{1}-x_{2}+S_{3}=-12$
and $x_{1}, x_{2}, S_{1}, S_{2}, S_{3} \geq 0$

| Iteration-1 |  | $C_{j}$ | -5 | -8 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $S_{1}$ | $S_{2}$ | $S_{3}$ |
| $S_{1}$ | 0 | -4 | -2 | -3 | 1 | 0 | 0 |
| $S_{2}$ | 0 | -5 | -1 | 2 | 0 | 1 | 0 |
| $S_{3}$ | 0 | -12 | (-1) | -1 | 0 | 0 | 1 |
| $Z=0$ |  | $Z_{j}$ | 0 | 0 | 0 | 0 | 0 |
|  |  | $C_{j}-Z_{j}$ | -5 | -8 | 0 | 0 | 0 |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{3}, j} \\ \text { and } S_{3}, j<0 \end{array}$ | $5 \uparrow$ | 8 | --- | --- | --- |

Minimum negative $X_{B}$ is -12 and its row index is 3 . So, the leaving basis variable is $S_{3}$.
Minimum positive ratio is 5 and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -1 .
Entering $=x_{1}$, Departing $=S_{3}$, Key Element $=-1$
$R_{3}($ new $)=R_{3}($ old $) \div-1$
$R_{1}$ (new) $=R_{1}$ (old) $+2 R_{3}$ (new)
$R_{2}($ new $)=R_{2}($ old $)+R_{3}$ (new)

| Iteration-2 |  | $C_{j}$ | -5 | -8 | 0 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{3}}$ |
| $S_{1}$ | 0 | 20 | 0 | -1 | 1 | 0 | -2 |
| $S_{2}$ | 0 | 7 | 0 | 3 | 0 | 1 | -1 |
| $x_{1}$ | -5 | 12 | 1 | 1 | 0 | 0 | -1 |
| $\boldsymbol{Z}=-\mathbf{6 0}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $-\mathbf{5}$ | $\mathbf{- 5}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{5}$ |
|  | $C_{j}-Z_{j}$ | 0 | -3 | 0 | 0 | -5 |  |
|  | Ratio | --- | --- | --- | --- | --- |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=12, x_{2}=0$
$\operatorname{Max} Z=-60$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN Z $=10 \times 1+20 \times 2$
subject to
$\mathrm{x} 1+2 \times 2>=6$
$\mathrm{x} 1+4 \times 2>=8$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=10 x_{1}+20 x_{2}$
subject to

$$
x_{1}+2 x_{2} \geq 6
$$

$$
x_{1}+4 x_{2} \geq 8
$$

and $x_{1}, x_{2} \geq 0$;

In order to apply the dual simplex method, convert Min $Z$ to Max $Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-10 x_{1}-20 x_{2}$
subject to

- $x_{1}-2 x_{2} \leq-6$
- $x_{1}-4 x_{2} \leq-8$
and $x_{1}, x_{2} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-10 x_{1}-20 x_{2}+0 S_{1}+0 S_{2}$
subject to

- $x_{1}-2 x_{2}+S_{1}=-6$
- $x_{1}-4 x_{2}+S_{2}=-8$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -10 | -20 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\boldsymbol{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | -6 | -1 | -2 | 1 | 0 |
| about:blank |  |  |  |  |  |  |


| $\boldsymbol{S}_{\mathbf{2}}$ | 0 | -8 | -1 | $\mathbf{( - 4 )}$ | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $C_{j}-Z_{j}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  |  | -10 | -20 | 0 | 0 |  |
|  | Ratio $=\frac{C_{j}-Z_{j}}{S_{2}, j}$ <br> and $S_{2}, j<0$ | 10 | $5 \uparrow$ | --- | --- |  |

Minimum negative $X_{B}$ is -8 and its row index is 2 . So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is 5 and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is -4 .
Entering $=x_{2}$, Departing $=S_{2}$, Key Element $=-4$
$R_{2}$ (new) $=R_{2}($ old $) \div-4$
$R_{1}$ (new) $=R_{1}$ (old) $+2 R_{2}$ (new)

| Iteration-2 |  | $C_{j}$ | -10 | -20 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $S_{1}$ | $S_{2}$ |
| $S_{1}$ | 0 | -2 | $\left(-\frac{1}{2}\right)$ | 0 | 1 | $-\frac{1}{2}$ |
| $x_{2}$ | -20 | 2 | $\frac{1}{4}$ | 1 | 0 | - $\frac{1}{4}$ |
| $Z=-40$ |  | $Z_{j}$ | -5 | -20 | 0 | 5 |
|  |  | $C_{j}-Z_{j}$ | -5 | 0 | 0 | -5 |
|  |  | $\begin{gathered} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{1}, j} \\ \text { and } S_{1}, j<0 \end{gathered}$ | $10 \uparrow$ | --- | --- | 10 |

Minimum negative $X_{B}$ is -2 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is 10 and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is $-\frac{1}{2}$.

Entering $=x_{1}$, Departing $=S_{1}$, Key Element $=-\frac{1}{2}$
$R_{1}($ new $)=R_{1}($ old $) \times-2$
$R_{2}($ new $)=R_{2}($ old $)-\frac{1}{4} R_{1}$ (new)

| Iteration-3 |  | $C_{j}$ | -10 | -20 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $x_{1}$ | -10 | 4 | 1 | 0 | -2 | 1 |
| $x_{2}$ | -20 | 1 | 0 | 1 | $\frac{1}{2}$ | $-\frac{1}{2}$ |
| $\boldsymbol{Z}=\mathbf{- 6 0}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $\mathbf{- 1 0}$ | $\mathbf{- 2 0}$ | $\mathbf{1 0}$ | $\mathbf{0}$ |
|  | $C_{j}-Z_{j}$ | 0 | 0 | -10 | 0 |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=4, x_{2}=1$
$\operatorname{Max} Z=-60$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=12 \times 1+8 \times 2$
subject to
$2 \times 1+2 \times 2>=6$
$3 \times 1+\mathrm{x} 2>=7$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=12 x_{1}+8 x_{2}$
subject to

$$
\begin{aligned}
& 2 x_{1}+2 x_{2} \geq 6 \\
& 3 x_{1}+x_{2} \geq 7 \\
& \text { and } x_{1}, x_{2} \geq 0
\end{aligned}
$$

In order to apply the dual simplex method, convert $\operatorname{Min} \mathrm{Z}$ to $\operatorname{Max} \mathrm{Z}$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-12 x_{1}-8 x_{2}$
subject to
$-2 x_{1}-2 x_{2} \leq-6$
$-3 x_{1}-x_{2} \leq-7$
and $x_{1}, x_{2} \geq 0 ;$

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type $' \leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-12 x_{1}-8 x_{2}+0 S_{1}+0 S_{2}$
subject to
$-2 x_{1}-2 x_{2}+S_{1}=-6$
$-3 x_{1}-x_{2}+S_{2}=-7$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -12 | -8 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | -6 | -2 | -2 | 1 | 0 |
|  |  |  |  |  |  |  |
| about:blank |  |  |  |  |  |  |


| $\boldsymbol{S}_{\mathbf{2}}$ | 0 | $\mathbf{- 7}$ | $\mathbf{( - 3 )}$ | -1 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $C_{j}-Z_{j}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  |  | Ratio $=\frac{C_{j}-Z_{j}}{S_{2}, j}$ <br> and $S_{2}, j<0$ | $4 \uparrow$ | -8 | 0 | 0 |
|  |  | 8 | --- | --- |  |  |

Minimum negative $X_{B}$ is -7 and its row index is 2 . So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is 4 and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -3 .
Entering $=x_{1}$, Departing $=S_{2}$, Key Element $=-3$
$R_{2}$ (new) $=R_{2}($ old $) \div-3$
$R_{1}$ (new) $=R_{1}$ (old) $+2 R_{2}$ (new)

| Iteration-2 |  | $C_{j}$ | -12 | -8 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $S_{1}$ | $S_{2}$ |
| $S_{1}$ | 0 | - $-\frac{4}{3}$ | 0 | $\left(-\frac{4}{3}\right)$ | 1 | $-\frac{2}{3}$ |
| $x_{1}$ | -12 | $\frac{7}{3}$ | 1 | $\frac{1}{3}$ | 0 | $-\frac{1}{3}$ |
| $Z=-28$ |  | $Z_{j}$ | -12 | -4 | 0 | 4 |
|  |  | $C_{j}-Z_{j}$ | 0 | -4 | 0 | -4 |
|  |  | $\begin{gathered} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{1}, j} \\ \text { and } S_{1}, j<0 \end{gathered}$ | --- | $3 \uparrow$ | --- | 6 |

Minimum negative $X_{B}$ is $-\frac{4}{3}$ and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is 3 and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is $-\frac{4}{3}$.

Entering $=x_{2}$, Departing $=S_{1}$, Key Element $=-\frac{4}{3}$
$R_{1}($ new $)=R_{1}($ old $) \times-\frac{3}{4}$
$R_{2}($ new $)=R_{2}($ old $)-\frac{1}{3} R_{1}$ (new)

| Iteration-3 |  | $C_{j}$ | -12 | -8 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $x_{2}$ | -8 | 1 | 0 | 1 | $-\frac{3}{4}$ | $\frac{1}{2}$ |
| $x_{1}$ | -12 | 2 | 1 | 0 | $\frac{1}{4}$ | $-\frac{1}{2}$ |
| $\boldsymbol{Z}=-\mathbf{3 2}$ |  | $Z_{\boldsymbol{j}}$ | $\mathbf{- 1 2}$ | $\mathbf{- 8}$ | $\mathbf{3}$ | $\mathbf{2}$ |
|  | $C_{j}-Z_{j}$ | 0 | 0 | -3 | -2 |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=2, x_{2}=1$
$\operatorname{Max} Z=-32$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=x 1+2 x 2$
subject to
$2 \times 1+\mathrm{x} 2>=4$
$\mathrm{x} 1+2 \times 2<=7$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=x_{1}+2 x_{2}$
subject to

$$
\begin{aligned}
2 x_{1}+x_{2} & \geq 4 \\
x_{1}+2 x_{2} & \leq 7
\end{aligned}
$$

and $x_{1}, x_{2} \geq 0$;

In order to apply the dual simplex method, convert Min $Z$ to Max $Z$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-x_{1}-2 x_{2}$
subject to
$-2 x_{1}-x_{2} \leq-4$

$$
x_{1}+2 x_{2} \leq 7
$$

and $x_{1}, x_{2} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-x_{1}-2 x_{2}+0 S_{1}+0 S_{2}$
subject to

$$
\begin{aligned}
-2 x_{1}-x_{2}+S_{1} & =-4 \\
x_{1}+2 x_{2}+S_{2} & =7
\end{aligned}
$$

and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -1 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $\boldsymbol{S}_{\mathbf{1}}$ | 0 | -4 | $\mathbf{( - 2 )}$ | -1 | 1 | 0 |
|  |  |  |  |  |  |  |
| about:blank |  |  |  |  |  |  |


| $S_{2}$ | 0 | 7 | 1 | 2 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $Z_{\boldsymbol{j}}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  |  | $C_{j}-Z_{j}$ | -1 | -2 | 0 | 0 |
|  |  | Ratio $=\frac{C_{j}-Z_{j}}{S_{1}, j}$ <br> and $S_{1}, j<0$ | $\frac{1}{2} \uparrow$ | 2 | --- | --e |

Minimum negative $X_{B}$ is -4 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is $\frac{1}{2}$ and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -2 .
Entering $=x_{1}$, Departing $=S_{1}$, Key Element $=-2$
$R_{1}($ new $)=R_{1}($ old $) \div-2$
$R_{2}$ (new) $=R_{2}$ (old) $-R_{1}$ (new)

| Iteration-2 |  | $C_{j}$ | -1 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $x_{1}$ | -1 | 2 | 1 | $\frac{1}{2}$ | $-\frac{1}{2}$ | 0 |
| $S_{2}$ | 0 | 5 | 0 | $\frac{3}{2}$ | $\frac{1}{2}$ | 1 |
| $\boldsymbol{Z}=-\mathbf{2}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $-\mathbf{1}$ | $-\frac{1}{2}$ | $\frac{\mathbf{1}}{\mathbf{2}}$ | $\mathbf{0}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | $-\frac{3}{2}$ | $-\frac{1}{2}$ | 0 |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=2, x_{2}=0$
$\operatorname{Max} Z=-2$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=\times 1+2 \times 2+2 \times 3$
subject to
$\mathrm{x} 1+\mathrm{x} 2+2 \times 3>=12$
$\mathrm{x} 1+2 \times 2+4 \times 3>=14$
and $\mathrm{x} 1, \mathrm{x} 2, \mathrm{x} 3>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=x_{1}+2 x_{2}+2 x_{3}$
subject to

$$
\begin{aligned}
& x_{1}+x_{2}+2 x_{3} \geq 12 \\
& x_{1}+2 x_{2}+4 x_{3} \geq 14
\end{aligned}
$$

and $x_{1}, x_{2}, x_{3} \geq 0$;

In order to apply the dual simplex method, convert $\operatorname{Min} \mathrm{Z}$ to $\operatorname{Max} \mathrm{Z}$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-x_{1}-2 x_{2}-2 x_{3}$
subject to

- $x_{1}-x_{2}-2 x_{3} \leq-12$
- $x_{1}-2 x_{2}-4 x_{3} \leq-14$
and $x_{1}, x_{2}, x_{3} \geq 0$;

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type ' $\leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-x_{1}-2 x_{2}-2 x_{3}+0 S_{1}+0 S_{2}$
subject to

- $x_{1}-x_{2}-2 x_{3}+S_{1}=-12$
$-x_{1}-2 x_{2}-4 x_{3}+S_{2}=-14$
and $x_{1}, x_{2}, x_{3}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -1 | -2 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{x}_{\mathbf{3}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | -12 | -1 | -1 | -2 | 1 | 0 |


| $S_{2}$ | 0 | -14 | -1 | -2 | $\mathbf{( - 4 )}$ | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $Z=\mathbf{0}$ |  | $Z_{j}$ | $C_{j}-Z_{j}$ | -1 | -2 | -2 | 0 |
|  |  | $C_{j}-Z_{j}$ <br> Ratio $=\frac{S_{2}, j}{}$ <br> and $S_{2}, j<0$ | 1 | 1 | $\frac{1}{2} \uparrow$ | --- | $\mathbf{0}$ |
|  |  | --- |  |  |  |  |  |

Minimum negative $X_{B}$ is -14 and its row index is 2 . So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is $\frac{1}{2}$ and its column index is 3 . So, the entering variable is $x_{3}$.
$\therefore$ The pivot element is -4 .
Entering $=x_{3}$, Departing $=S_{2}$, Key Element $=-4$
$R_{2}($ new $)=R_{2}($ old $) \div-4$
$R_{1}($ new $)=R_{1}($ old $)+2 R_{2}($ new $)$

| Iteration-2 |  | $C_{j}$ | -1 | -2 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $S_{1}$ | $S_{2}$ |
| $S_{1}$ | 0 | -5 | $\left(-\frac{1}{2}\right)$ | 0 | 0 | 1 | $-\frac{1}{2}$ |
| $x_{3}$ | -2 | $\frac{7}{2}$ | $\frac{1}{4}$ | $\frac{1}{2}$ | 1 | 0 | - $\frac{1}{4}$ |
| $Z=-7$ |  | $Z_{j}$ | - $\frac{1}{2}$ | -1 | -2 | 0 | $\frac{1}{2}$ |
|  |  | $C_{j}-Z_{j}$ | - $\frac{1}{2}$ | -1 | 0 | 0 | - $\frac{1}{2}$ |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{1}, j} \\ \text { and } S_{1}, j<0 \end{array}$ | $1 \uparrow$ | --- | --- | --- | 1 |

Minimum negative $X_{B}$ is -5 and its row index is 1 . So, the leaving basis variable is $S_{1}$.

Minimum positive ratio is 1 and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is $-\frac{1}{2}$.

Entering $=x_{1}$, Departing $=S_{1}$, Key Element $=-\frac{1}{2}$
$R_{1}($ new $)=R_{1}($ old $) \times-2$
$R_{2}($ new $)=R_{2}($ old $)-\frac{1}{4} R_{1}($ new $)$

| Iteration-3 |  | $C_{j}$ | -1 | -2 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{x}_{\mathbf{3}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $x_{1}$ | -1 | 10 | 1 | 0 | 0 | -2 | 1 |
| $x_{3}$ | -2 | 1 | 0 | $\frac{1}{2}$ | 1 | $\frac{1}{2}$ | $-\frac{1}{2}$ |
| $\boldsymbol{Z}=-\mathbf{1 2}$ |  | $\boldsymbol{Z}_{\boldsymbol{j}}$ | $-\mathbf{1}$ | $\mathbf{- 1}$ | $-\mathbf{2}$ | $\mathbf{1}$ | $\mathbf{0}$ |
|  | $C_{j}-Z_{j}$ | 0 | -1 | 0 | -1 | 0 |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.

Hence, optimal solution is arrived with value of variables as :
$x_{1}=10, x_{2}=0, x_{3}=1$
$\operatorname{Max} Z=-12$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=\mathrm{x} 1+\mathbf{2 x} 2$
subject to
$-2 \times 1-\times 2<=-4$
$-\mathrm{x} 1-2 \times 2<=-7$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=x_{1}+2 x_{2}$
subject to
$-2 x_{1}-x_{2} \leq-4$
$-x_{1}-2 x_{2} \leq-7$
and $x_{1}, x_{2} \geq 0$;

In order to apply the dual simplex method, convert Min Z to Max Z
Problem is
$\operatorname{Max} Z=-x_{1}-2 x_{2}$
subject to
$-2 x_{1}-x_{2} \leq-4$
$-x_{1}-2 x_{2} \leq-7$
and $x_{1}, x_{2} \geq 0 ;$

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type $' \leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-x_{1}-2 x_{2}+0 S_{1}+0 S_{2}$
subject to
$-2 x_{1}-x_{2}+S_{1}=-4$
$-x_{1}-2 x_{2}+S_{2}=-7$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -1 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | -4 | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ |
| $S_{1}$ | 0 | -7 | -2 | -1 | 1 | 0 |
|  | 0 | $(-\mathbf{1})$ | -2 | 0 | 1 |  |


| $\boldsymbol{S}_{\mathbf{2}}$ |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $Z_{\boldsymbol{j}}$ | $C_{j}-Z_{j}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  |  | -1 | -2 | 0 | 0 |  |
|  | Ratio $=\frac{C_{j}-Z_{j}}{S_{2}, j}$ <br> and $S_{2}, j<0$ | $1 \uparrow \uparrow$ | 1 | --- | --- |  |

Minimum negative $X_{B}$ is -7 and its row index is 2 . So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is 1 and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -1 .
Entering $=x_{1}$, Departing $=S_{2}$, Key Element $=-1$
$R_{2}$ (new) $=R_{2}($ old $) \div-1$
$R_{1}$ (new) $=R_{1}$ (old) $+2 R_{2}$ (new)

| Iteration-2 |  | $C_{j}$ | -1 | -2 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | 10 | 0 | 3 | 1 | -2 |
| $x_{1}$ | -1 | 7 | 1 | 2 | 0 | -1 |
| $\boldsymbol{Z}=-7$ |  | $Z_{\boldsymbol{j}}$ | $-\mathbf{1}$ | $-\mathbf{2}$ | $\mathbf{0}$ | $\mathbf{1}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | 0 | 0 | -1 |
|  | Ratio | --- | --- | $-\boldsymbol{-}$ | $\boldsymbol{- - -}$ |  |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=7, x_{2}=0$
$\operatorname{Max} Z=-7$

Solution is provided by AtoZmath.com

## Print This Solution Close This Solution

Find solution using dual-simplex method
MIN $Z=x 1+x 2$
subject to
$\mathrm{x} 1+3 \times 2>=6$
$2 \times 1+\mathrm{x} 2>=8$
and $\mathrm{x} 1, \mathrm{x} 2>=0$

## Solution:

## Problem is

$\operatorname{Min} Z=x_{1}+x_{2}$
subject to

$$
\begin{array}{r}
x_{1}+3 x_{2} \geq 6 \\
2 x_{1}+x_{2} \geq 8 \\
\text { and } x_{1}, x_{2} \geq 0
\end{array}
$$

In order to apply the dual simplex method, convert $\operatorname{Min} \mathrm{Z}$ to $\operatorname{Max} \mathrm{Z}$ and all $\geq$ constraint to $\leq$ constraint by multiply -1 .

## Problem is

$\operatorname{Max} Z=-x_{1}-x_{2}$
subject to

- $x_{1}-3 x_{2} \leq-6$
$-2 x_{1}-x_{2} \leq-8$
and $x_{1}, x_{2} \geq 0 ;$

The problem is converted to canonical form by adding slack, surplus and artificial variables as appropiate

1. As the constraint 1 is of type ' $\leq$ ' we should add slack variable $S_{1}$
2. As the constraint 2 is of type $' \leq$ ' we should add slack variable $S_{2}$

## After introducing slack variables

$\operatorname{Max} Z=-x_{1}-x_{2}+0 S_{1}+0 S_{2}$
subject to
$-x_{1}-3 x_{2}+S_{1}=-6$
$-2 x_{1}-x_{2}+S_{2}=-8$
and $x_{1}, x_{2}, S_{1}, S_{2} \geq 0$

| Iteration-1 |  | $C_{j}$ | -1 | -1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $S_{1}$ | 0 | -6 | -1 | -3 | 1 | 0 |
|  |  |  |  |  |  |  |
| about:blank |  |  |  |  |  |  |


| $\boldsymbol{S}_{\mathbf{2}}$ | 0 | -8 | $\mathbf{( - 2 )}$ | -1 | 0 | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{Z}=\mathbf{0}$ |  | $Z_{j}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
|  |  | $C_{j}-Z_{j}$ | -1 | -1 | 0 | 0 |
|  | Ratio $=\frac{C_{j}-Z_{j}}{S_{2}, j}$ <br> and $S_{2}, j<0$ | $\frac{1}{2} \uparrow$ | 1 | --- | --- |  |
|  |  |  |  |  |  |  |

Minimum negative $X_{B}$ is -8 and its row index is 2 . So, the leaving basis variable is $S_{2}$.
Minimum positive ratio is $\frac{1}{2}$ and its column index is 1 . So, the entering variable is $x_{1}$.
$\therefore$ The pivot element is -2 .
Entering $=x_{1}$, Departing $=S_{2}$, Key Element $=-2$
$R_{2}$ (new) $=R_{2}($ old $) \div-2$
$R_{1}$ (new) $=R_{1}$ (old) $+R_{2}$ (new)

| Iteration-2 |  | $C_{j}$ | -1 | -1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | $C_{B}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $x_{1}$ | $x_{2}$ | $S_{1}$ | $S_{2}$ |
| $S_{1}$ | 0 | -2 | 0 | $\left(-\frac{5}{2}\right)$ | 1 | $-\frac{1}{2}$ |
| $x_{1}$ | -1 | 4 | 1 | $\frac{1}{2}$ | 0 | $-\frac{1}{2}$ |
| $Z=-4$ |  | $Z_{j}$ | -1 | $-\frac{1}{2}$ | 0 | $\frac{1}{2}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | - $\frac{1}{2}$ | 0 | $-\frac{1}{2}$ |
|  |  | $\begin{array}{r} \text { Ratio }=\frac{C_{j}-Z_{j}}{S_{1}, j} \\ \text { and } S_{1}, j<0 \end{array}$ | --- | $\frac{1}{5} \uparrow$ | --- | 1 |

Minimum negative $X_{B}$ is -2 and its row index is 1 . So, the leaving basis variable is $S_{1}$.
Minimum positive ratio is $\frac{1}{5}$ and its column index is 2 . So, the entering variable is $x_{2}$.
$\therefore$ The pivot element is $-\frac{5}{2}$.

Entering $=x_{2}$, Departing $=S_{1}$, Key Element $=-\frac{5}{2}$
$R_{1}($ new $)=R_{1}($ old $) \times-\frac{2}{5}$
$R_{2}($ new $)=R_{2}($ old $)-\frac{1}{2} R_{1}($ new $)$

| Iteration-3 |  | $C_{j}$ | -1 | -1 | 0 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{B}$ | $\boldsymbol{C}_{\boldsymbol{B}}$ | $\boldsymbol{X}_{\boldsymbol{B}}$ | $\boldsymbol{x}_{\mathbf{1}}$ | $\boldsymbol{x}_{\mathbf{2}}$ | $\boldsymbol{S}_{\mathbf{1}}$ | $\boldsymbol{S}_{\mathbf{2}}$ |
| $x_{2}$ | -1 | $\frac{4}{5}$ | 0 | 1 | $-\frac{2}{5}$ | $\frac{1}{5}$ |
| $x_{1}$ | -1 | $\frac{18}{5}$ | 1 | 0 | $\frac{1}{5}$ | $-\frac{3}{5}$ |
| $\boldsymbol{Z}=-\frac{\mathbf{2 2}}{\mathbf{5}}$ |  | $Z_{j}$ | $-\mathbf{1}$ | $-\mathbf{1}$ | $\frac{\mathbf{1}}{\mathbf{5}}$ | $\frac{\mathbf{2}}{\mathbf{5}}$ |
|  |  | $C_{j}-Z_{j}$ | 0 | 0 | $-\frac{1}{5}$ | $-\frac{2}{5}$ |

Since all $C_{j}-Z_{j} \leq 0$ and all $X_{B i} \geq 0$ thus the current solution is the optimal solution.
Hence, optimal solution is arrived with value of variables as :
$x_{1}=\frac{18}{5}, x_{2}=\frac{4}{5}$
$\operatorname{Max} Z=-\frac{22}{5}$

Solution is provided by AtoZmath.com

